

FA2022 Week 07

# Crypto I

Anakin



# Outline

Basics

XOR

Diffie-Hellman



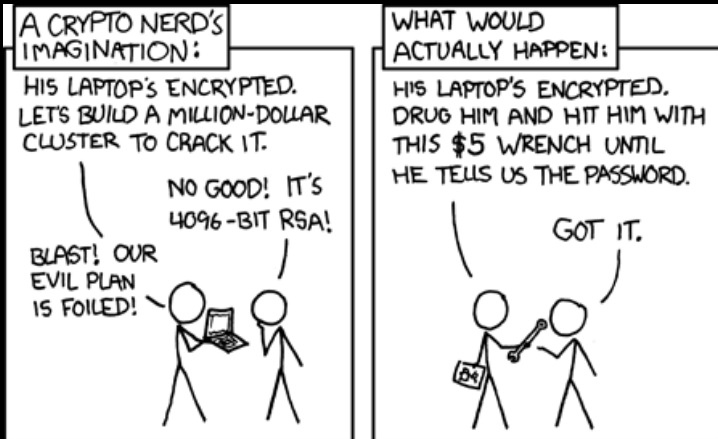
# Announcements

- ACM Cleanup Sunday after Crypto II



sigpwny{n0t\_that\_crypt0\_but\_th3\_0th3r\_0n3}

ctf.sigpwny.com



# Section 1

## Basics



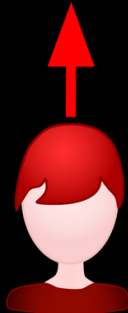
# What is Crypto Anyways?



Alice



Bob



Eve



# Why Do We Care?

https://



> \_SSH



# Crypto in Ye Olden Days

- Relied on simple patterns
- Hard / annoying to break by hand, **easy to break by computer**
- Examples:
  - Caesar Cipher (rot k)
    - $a \rightarrow c, b \rightarrow d, \dots, y \rightarrow a, z \rightarrow b$  (rot 2)
  - Substitution
    - Create a table mapping each letter to another
    - Generalization of Caesar Cipher
  - Many More
    - **All insecure!!**





# Data Representation

- TL;DR: computers store things in binary (0s and 1s), and we have different ways of representing this
- Tip: In Python, always work with bytes / bytestrings, never with normal strings (Python 3.8+)
- Look at the challenge source if given and mimic what they do



# Conversion Cheatsheet

This is hard to read, download the slides!!

Format	Description	From Bytes	To Bytes
<b>base64</b>	uses printable letters to encode more complex binary	<code>base64.b64encode</code>	<code>base64.b64decode</code>
hex	uses symbols 0-9, A-F	<code>bytearray.hex()</code> , <code>binascii.hexlify()</code>	<code>bytes.fromhex()</code> , <code>binascii.unhexlify()</code>
integer	normal integers	<code>Crypto.Util.number.bytes_to_long</code> (PyCryptoDome), <code>int.from_bytes</code>	<code>Crypto.Util.number.long_to_bytes</code> (PyCryptoDome), <code>int.to_bytes</code>

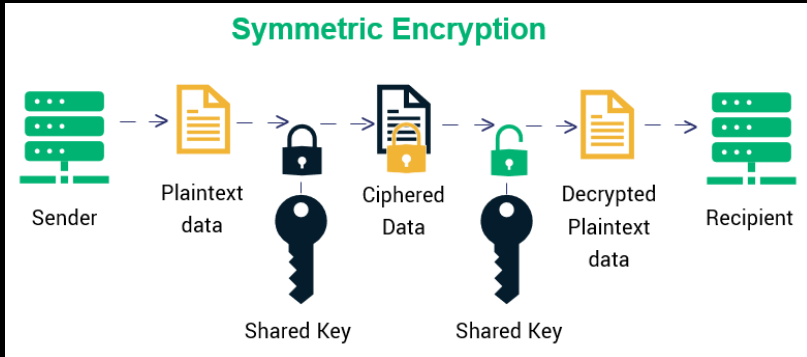


## Section 2

**XOR**



# Symmetric Encryption



# XOR

A	B	$A \oplus B$
0	0	0
0	1	1
1	0	1
1	1	0



# XOR

- XOR has some really nice properties that make it perfect for symmetric encryption
- Say  $M$  is some message as a bitstring,  $K$  is some key
- Then let  $C = M \oplus K$  be a ciphertext
- Properties:
  - **Order doesn't matter:**  $M \oplus K = K \oplus M$
  - **Group as needed:**  $M \oplus (K \oplus K) = (M \oplus K) \oplus K$
  - **$\emptyset$  is the identity:**  $M \oplus \emptyset = M$
  - **Self Inverse:**  $K \oplus K = \emptyset$
- All of this means  $C \oplus K = M \oplus K \oplus K = M \oplus \emptyset = M$



# Overview of Some Attacks

- For certain reasons, in general XOR is really really hard to break
  - Without more information, you need to try  $2^\lambda$  guesses to break a bitstring of length  $\lambda$
- Usually you need to know some information about the plaintext
  - Known plaintext
  - Properties like language
- You may need to know some information about the key
  - Really short keys are able to be brute forced
  - Check if code leaks information about key length



## Section 3

# Diffie-Hellman





# Modular Arithmetic

- In cryptography, numbers can get really big really fast
- We use **modular arithmetic** to deal with this
- Modular arithmetic is **arithmetic with remainders after division**



# Remainders

- Assume we have some number  $n$ . We are going to do some computation **mod**  $n$
- For now, say  $n = 101$

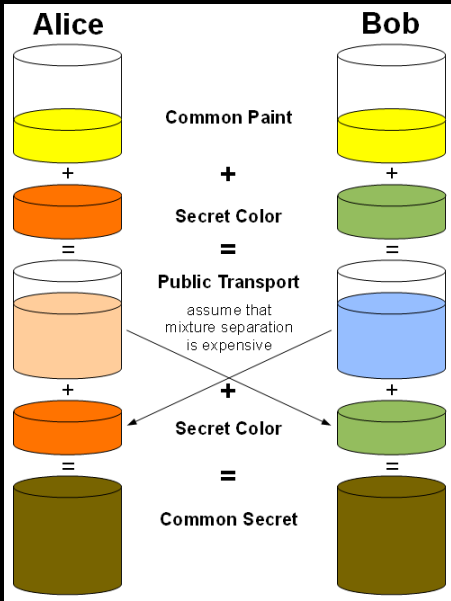
$$\begin{aligned}131 + 140 * (102)^{2000} &\equiv 131 + 39 * (102)^{2000} && (\text{mod } 101) \\ &\equiv 30 + 39 * (102)^{2000} && (\text{mod } 101) \\ &\equiv 30 + 39 * (1)^{2000} && (\text{mod } 101) \\ &\equiv 30 + 39 && (\text{mod } 101) \\ &\equiv 69 && (\text{mod } 101)\end{aligned}$$



# Discrete Log

- If  $a^b \equiv X \pmod{p}$ ,  $b$  is the **discrete log of  $X$  with base  $a$** .
- Given some random  $X$  and  $a$ , finding  $b$  is really hard to compute for large primes  $p$
- This **Discrete Log Problem** is the basis for many modern cryptography standards





# Painting with Numbers

- Let  $g$  be a public number we call a generator and  $p$  be some public prime
- Alice generates secret  $a$  and computes  $A \equiv g^a \pmod{p}$
- Bob generates secret  $b$  and computes  $B \equiv g^b \pmod{p}$
- Alice sends Bob  $A$  and Bob sends Alice  $B$
- Alice computes  $B^a \pmod{p}$
- Bob computes  $A^b \pmod{p}$

$$A^b \equiv (g^a)^b \equiv g^{ab} \equiv (g^b)^a \equiv B^a \pmod{p}$$



# Overview of Some Attacks

Remember, discrete log in general is **hard**

- Small Primes
- Primes are generated in specific ways (“smooth primes”)
- “Oracle” attacks (access to a special machine that leaks information)

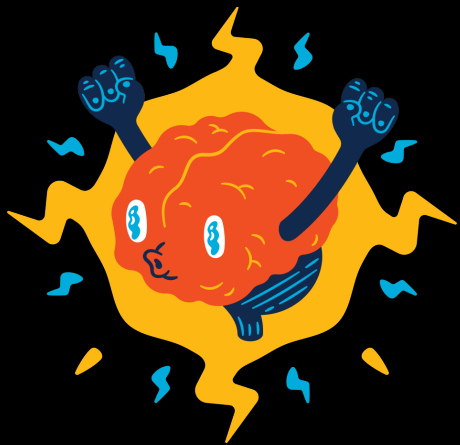


# Tools!

- Python + [SageMath](#) is your friend
- [PyCryptodome](#) is an extremely useful Python crypto library
- [PwnTools](#) will allow you to automate parts of your attacks
- Google + StackOverflow (“how to crack DH with ...”)
- Installation is **annoying**, use the [CryptoHack Docker](#)



# Practice @ CryptoHack





# Next Meetings

## 2022-10-16 – This Sunday

- Crypto II
- More Diffie-Hellman + RSA
- ACM Cleanup afterwards

## 2022-10-20 – Next Thursday

- Rev II with Richard
- angr + Z3

## 2022-10-23 – Next Sunday

- Research Presentation from Mingjia
- Stealing Hospital Information





**SIGPwny**